

Compressive Sensing in Clifford Analysis: Bicomplex and Quaternionic case

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Overview

- 1 Compressive Sensing (CS)
 - What is it?
 - The Mathematical Setting
 - Restricted Isometry Property (RIP)
- 2 Bicomplex Setting
 - Fourier Basis and Setting in \mathbb{R}^4 , \mathbb{BC} -DFT
 - CS for bicomplex signal
- 3 Quaternionic Setting
 - Fourier Basis and Setting in \mathbb{R}^4 , \mathbb{H} -DFT
 - CS for quaternionic signal
- 4 Results and further challenges

What is it?

This field [Compressive Sensing or Compressed Sensing \(CS\)](#) started with [Candes/Romberg/Tao 2006](#) which showed that the number of measurements can be small and still contain nearly all the useful information required for an [effective reconstruction](#).

Main requirements

- a signal is expressible in terms of a linear combination of elements in a given basis ([frame](#))
- it has a sparse representation

What is the main problem?

Suppose there is a signal $\mathbf{x} = (x_1, x_2, \dots, x_N)^T$.

From this we measure just a few linear samples $\mathbf{y} = (y_1, y_2, \dots, y_M)^T$, with $M \ll N$.

Can we reconstruct \mathbf{x} ?

The answer is **yes** if \mathbf{x} is sparse in a given basis or frame.

Graphic representation

Solve $y = \Phi x$ such that

$$\begin{array}{c}
 y \\
 \begin{array}{|c|} \hline \text{colored squares} \\ \hline \end{array} \\
 M \times 1
 \end{array}
 =
 \begin{array}{c}
 \Phi \\
 \begin{array}{|c|} \hline \text{colored grid} \\ \hline \end{array} \\
 M \times N \ (M < N)
 \end{array}
 \begin{array}{c}
 x \\
 \begin{array}{|c|} \hline \text{white squares with some colored} \\ \hline \end{array} \\
 N \times 1
 \end{array}$$

Restricted Isometry Property - RIP

Definition

Φ has the **RIP** property if for each integer $k \in \mathbb{N}$ there exists an isometry constant $0 < \delta_k < 1$ that is, δ_k is the smallest positive real number such that it holds

$$(1 - \delta_k) \|x\|_{\ell_2}^2 \leq \|\Phi x\|_{\ell_2}^2 \leq (1 + \delta_k) \|x\|_{\ell_2}^2$$

for all k -sparse vectors x .

When is x reconstructible?

- if the signal x is k -sparse, then only $M \geq 2k$ data points are required for reconstruction; moreover,
- stability is attained if $M \geq Ck \ln(N/M)$.

Big Problem

RIP property is a too **strong** condition to demand.

SOLUTION: find a condition which can be fulfilled with a certain probability such that we have reconstruction of a signal by an

- ℓ_1 -minimization procedure;
- uniformly distributed randomly chosen sampling points.

Fundamental Lemma

Let $c \in \ell_2(D)$ and $T := \text{supp} c$. Assume

$$\mathcal{F}_{TX} : \ell_2(T) \rightarrow \ell_2(X)$$

to be injective. Suppose that there exists a vector $P \in \ell_2(D)$ with the following properties:

- (i) $P_k = \text{sgn} c_k$ for all $k \in T$;
- (ii) $|P_k| < 1$ for all $k \notin T$;
- (iii) there exists a vector $\lambda \in \ell_2(X)$ such that $P = \mathcal{F}_X^* \lambda$ then c is unique minimizer to our ℓ_1 -minimization problem

$$\min ||c_k||_1 := \sum_{k \in T} |c_k|,$$

s.t.

$$f(x_j) := \sum_{k \in T} c_k \varphi(x_j), \quad j = 1, \dots, N.$$

Which cases we already study?

Cases with results

- ① Takenaka-Malmquist System Setting
- ② Bicomplex Setting
- ③ Quaternionic Setting
- ④ Slice Monogenic Setting

Cases without results

- ① Spherical harmonics Setting

In this presentation:

- ① Bicomplex Setting
- ② Quaternionic Setting

Takenaka-Malmquist System Setting

- Compressed Sensing with Nonlinear Fourier Atoms, Narciso Gomes, Stefan Hartmann and Paula Cerejeiras, Springer Birkhäuser, Cham (2016)
- <https://link.springer.com/chapter/10.1007/978-3-319-42529-04>

The Takenaka-Malmquist system

The **Takenaka-Malmquist (TM) system** with respect to a sequence $\{a_0, a_1, \dots, a_n\}$ of points in the unit disk is given by

$$\mathcal{B}_0(z) = \frac{\sqrt{1 - |a_0|^2}}{1 - \overline{a_0}z},$$

$$\mathcal{B}_k(z) = \frac{\sqrt{1 - |a_k|^2}}{1 - \overline{a_k}z} \prod_{l=0}^{k-1} \frac{z - a_l}{1 - \overline{a_l}z}, \quad k = 1, 2, \dots$$

Slice regular functions

Definition

Let U be an open set in \mathbb{H} and let $f : U \rightarrow H$ be real differentiable. The function f is said to be (left) slice regular or (left) slice hyperholomorphic if for every $I \in \mathbb{S}$, its restriction f_I to the complex plane $\mathbb{C}_I = \mathbb{R} + I\mathbb{R}$ passing through origin and containing I and satisfies

$$\bar{\partial}_I f(x + Iy) := \frac{1}{2} \left(\frac{\partial}{\partial x} + I \frac{\partial}{\partial y} \right) f_I(x + Iy) = 0$$

on $U \cap \mathbb{C}_I$. The class of (left) slice regular functions on U . Analogously, a function is said to be right slice regular in U if

$$(f_I \bar{\partial}) f(x + Iy) := \frac{1}{2} \left(\frac{\partial}{\partial x} f_I(x + Iy) + \frac{\partial}{\partial y} f_I(x + Iy) I \right) = 0.$$

Representation Formula

Representation Formula

Given $s = x + Jy \in \mathbb{C}_J$, there are unique \mathbb{C}_J such that

$$f(x + Jy) = (1 - JI)f(x + Iy) + (1 + JI)f(x - Iy)$$

$$\Leftrightarrow$$

$$e^{Jk\theta} = \frac{1}{2} \left[(1 - JI)e^{Ik\theta}(1 + JI) \right].$$

Bicomplex Setting

Bicomplex Setting

- Bicomplex signals with sparsity constraints, Narciso Gomes, Paula Cerejeiras, Wiley Library, (2018)
- <https://onlinelibrary.wiley.com/doi/pdf/10.1002/mma.5059>

Fourier basis

For the description of an image one uses the standard Fourier basis:

$$\left\{ e^{i\theta}, \theta \in [0, 2\pi[\right\}. \quad (1)$$

The **discrete Fourier transform (DFT)** converts the sampled function from the time-domain to the frequency domain, that is to say, it converts a finite list of equally-spaced samples of a function into the list of coefficients for complex sinusoids.

For our setting we use the set of all automorphisms of the unit ball $\varphi : \mathbb{B}_1 \rightarrow \mathbb{B}_1$, (\mathbb{B}_1 the unit ball in \mathbb{BC}) restricted to the boundary. This set can be regarded as a generalization of the classical Fourier system (1).

Motivation

Bidimensional extension - image processing

- An approach for analyzing image is the concept of “bicomplex signal”.
- Embedding of a multi-channel signal into the bicomplex algebra.
- Representation of HSI-model
((hue-saturation-intensity)-model for human eye vision)
- The **idempotent** property allow us to use a similar approach to the one-dimensional case.

Bicomplex numbers

Definition

$$\mathbb{BC} := \{\zeta = \zeta_1 + \mathbf{j}\zeta_2 : \zeta_1 = a + \mathbf{i}b, \zeta_2 = c + \mathbf{i}d \in \mathbb{C}\},$$

where

$$\mathbf{j}^2 = \mathbf{i}^2 = -1 \quad \text{and} \quad \mathbf{ij} = \mathbf{ji}.$$

Important: the imaginary units \mathbf{i} and \mathbf{j} are **commuting**!

Problem: generalization to high dimension implies loss of certain properties!

Zero Divisors

Existence of **zero divisors**:

$$\mathbf{e}^+ := \frac{1 + \mathbf{i}\mathbf{j}}{2} \quad \text{and} \quad \mathbf{e}^- := \frac{1 - \mathbf{i}\mathbf{j}}{2}.$$

Properties:

- Unitary sum: $\mathbf{e}^+ + \mathbf{e}^- = 1$;
- Difference property: $\mathbf{e}^+ - \mathbf{e}^- = \mathbf{i}\mathbf{j}$;
- Idempotent property: $(\mathbf{e}^+)^2 = \mathbf{e}^+$, $(\mathbf{e}^-)^2 = \mathbf{e}^-$.

The property of **idempotent** allows to separate the space of bicomplex in terms of \mathbf{e}^+ , \mathbf{e}^- as projectors.

Decomposition

Decomposition

Given $\zeta = \zeta_1 + \mathbf{j}\zeta_2 \in \mathbb{BC}$, ($\zeta_1, \zeta_2 \in \mathbb{C}$), there are unique $z, w \in \mathbb{C}$ such that

$$\zeta = ze^+ + we^-,$$

with $z = \zeta_1 - \mathbf{i}\zeta_2$ and $w = \zeta_1 + \mathbf{i}\zeta_2$.

Conjugation

The correspondent **conjugation** of a given bicomplex number $\zeta = ze^+ + we^-$ is defined as

$$\bar{\zeta} = \bar{z}e^+ + \bar{w}e^-.$$

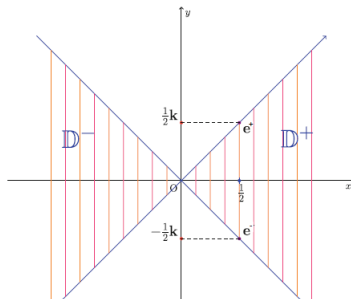
Hyperbolic positive numbers

The idempotent representation of a bicomplex number

$\zeta = a + b\mathbf{i}\mathbf{j} = a + b\mathbf{k}$, with $a, b \in \mathbb{R}$, is $\zeta = \nu\mathbf{e}^+ + \mu\mathbf{e}^-$ such that $\nu, \mu \in \mathbb{R}$, with $\nu = b + a$ and $\mu = b - a$.

Then, we define the set of positive hyperbolic numbers as

$$\mathbb{D}^+ := \{\nu\mathbf{e}^+ + \mu\mathbf{e}^- \mid \nu, \mu \geq 0\}.$$



Inner product and related norms

- The “inner product” of two bicomplex numbers, $\zeta = z_1 \mathbf{e}^+ + w_1 \mathbf{e}^-$ and $\eta = z_2 \mathbf{e}^+ + w_2 \mathbf{e}^-$, is defined as

$$\langle \zeta, \eta \rangle := z_1 \bar{z}_2 \mathbf{e}^+ + w_1 \bar{w}_2 \mathbf{e}^- = \langle z_1, z_2 \rangle_{\mathbb{C}} \mathbf{e}^+ + \langle w_1, w_2 \rangle_{\mathbb{C}} \mathbf{e}^-.$$

- The “bicomplex norm” is defined as

$$|\zeta|_{\mathbf{k}} := |z_1| \mathbf{e}^+ + |z_2| \mathbf{e}^- \in \mathbb{D}^+.$$

Then the induced (euclidean) norm by this inner product is

$$\|\zeta\| := 2\operatorname{Re}(|\zeta|_{\mathbf{k}}).$$

Inverse

An invertible bicomplex number $\zeta = ze^+ + we^-$ has its inverse given by

$$\zeta^{-1} = z^{-1}e^+ + w^{-1}e^-$$

where z^{-1} and w^{-1} are the complex multiplicative inverses of z and w , respectively.

High dimension extension

In \mathbb{R}^2 the classical discrete Fourier transform of a signal f is given by

$$F(x_1, x_2) = \sum_{m=t_1}^{M-1} \sum_{t_2=0}^{N-1} f(t_1, t_2) e^{i2\pi\left(\frac{t_1 x_1}{M} + \frac{t_2 x_2}{N}\right)}$$

Definition (\mathbb{BC} –Discrete Fourier Transform)

Given a signal f in \mathbb{R}^2 , we define its \mathbb{BC} –discrete Fourier transform as

$$F(x_j^1, x_j^2) = \sum_{k_1, k_2 \in T} f(k_1, k_2) \left[e^{ik_1 x_j^1} \mathbf{e}^+ + e^{ik_2 x_j^2} \mathbf{e}^- \right]$$

Our setting

Denote by Π_q the space of functions spanned by $d = (2q + 1)^2$ elements of the bicomplex system.

Additionally, we consider the signal f to have its support on a set T satisfying to $|T| \leq M \ll d$, that is to say,

$$f(x) = \sum_{(k, \tilde{k}) \in T} c_{k\tilde{k}} \left(e^{ix^1} \mathbf{e}^+ + e^{ix^2} \mathbf{e}^- \right),$$

with the sequence $(c_{k\tilde{k}} := c_k \mathbf{e}^+ + c_{\tilde{k}} \mathbf{e}^-)_{k\tilde{k}} \in \mathbb{BC}$ having support on T and $x = (x^1, x^2) \in \mathbb{R}^2$.

Main Goal

Given a sampling set

$$X := \{x_1, x_2, \dots, x_N\}$$

of linearly independent random variables (l.i.r.v.) having uniform distribution on $[0, 2\pi]^2$ our goal is to reconstruct the signal f from the sparse samples

$$\{f(x_j), j = 1, 2, \dots, N\},$$

and this with a certain (known) probability.

Main theorem

Theorem

Assume $f \in \prod_q$ with some sparsity $M \in \mathbb{N}$, and consider $X := \{x_1, x_2, \dots, x_N\}$ a sampling set of *i.i.v.* with uniform distribution on $[0, 2\pi]^2$.

Choose $n \in \mathbb{N}$, $\beta > 0$, $\kappa > 0$ and $K_1, K_2, \dots, K_n \in \mathbb{N}$ s. t.

$$a := \sum_{m=1}^n \beta^{n/K_m} < 1 \quad \text{and} \quad \frac{\kappa}{1-\kappa} \leq \frac{1-a}{1+a} M^{-3/2}.$$

For $\theta := N/M$ we have that f can be reconstructed with probability at least

$$1 - \left(D\beta^{-2n} \sum_{m=1}^n G_{2mK_m}(\theta) + M\kappa^{-2} G_{2n}(\theta) \right)$$

from its sample values $f(x_1), \dots, f(x_N)$ by solving the following ℓ^1 -minimization problem (2).

ℓ^1 -minimization problem

$$\min ||(c_{k\tilde{k}})||_1 := \sum_{k \in T} |c_{k\tilde{k}}|, \quad (2)$$

subject to

$$g(x_j) := \sum_{(k, \tilde{k}) \in T} c_{k\tilde{k}} \left(e^{ix_j^1} \mathbf{e}^+ + e^{ix_j^2} \mathbf{e}^- \right), \quad j = 1, \dots, N.$$

Hereby,

$$G_n(\theta) = \theta^{-n} \sum_{k=1}^{\lfloor n/2 \rfloor} S_2(n, k) \theta^k$$

where $S_2(n, k)$ denotes the Stirling numbers of second kind.

Proof - notations

Some **auxiliary notations**:

$\ell^2(D)$, $\ell^2(T)$, $\ell^2(X)$ will denote the ℓ^2 —spaces of sequences in those spaces endowed with the usual **Euclidean norm**.

The operator $\mathcal{F}_X : \ell^2(D) \rightarrow \ell^2(X)$.

We also consider the operator \mathcal{F}_{TX} as the **restriction** of \mathcal{F}_X which denote the sequences with **support** only on T acting from $\ell^2(T)$ in $\ell^2(X)$, and the their **adjoint** operators \mathcal{F}_{TX}^* ,

$$\mathcal{F}_X^* : \ell^2(X) \rightarrow \ell^2([-q, q]^2) \text{ and } \mathcal{F}_{TX}^* : \ell^2(X) \rightarrow \ell^2(T).$$

Sampling matrix

We consider the **sampling matrix**

$$\mathcal{F}_X := \left[e^{ikx_\ell^1} \mathbf{e}^+ + e^{i\tilde{k}x_\ell^2} \mathbf{e}^- \right]_{k, \tilde{k}=1, \dots, d, \ell=1, \dots, N}.$$

We define the **sign function** in bicomplex case as

$$\operatorname{sgn} c := \left(\frac{c_k}{|c_k|} \mathbf{e}^+ + \frac{c_{\tilde{k}}}{|c_{\tilde{k}}|} \mathbf{e}^- \right).$$

We consider the following **necessary lemmas**:

Proof - necessary lemma

Lemma

Let f and g be two bicomplex functions. It holds

- ① $|\operatorname{sgn} f|_{\mathbf{k}} = 1$;
- ② $f \overline{\operatorname{sgn} f} = |f|_{\mathbf{k}}$;
- ③ $|g|_{\mathbf{k}} |\operatorname{sgn} f|_{\mathbf{k}} = |g \operatorname{sgn} f|_{\mathbf{k}}$.

Lemma

For a non-singular matrix $\mathcal{F}_{TX}^* \mathcal{F}_{TX}$ we have

$$(\mathcal{F}_{TX}^* \mathcal{F}_{TX})^{-1} = [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_1^{-1} \mathbf{e}^+ + [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_2^{-1} \mathbf{e}^-.$$

Proof - necessary lemma

Lemma

Let $c_k \mathbf{e}^+ + c_{\tilde{k}} \mathbf{e}^- =: c \in \ell^2(D)$ and $T := \text{supp} c$. Assume

$$\mathcal{F}_{TX} : \ell^2(T) \rightarrow \ell^2(X)$$

to be injective. Suppose that there exists a vector $P \in \ell^2(D)$ with the following properties:

- (i) $P_{k\tilde{k}} = \text{sgn}(c_{k\tilde{k}})_{k\tilde{k}}$ for all $(k, \tilde{k}) \in T$,
- (ii) $|P_{k\tilde{k}}| < 1$ for all $(k, \tilde{k}) \notin T$,
- (iii) there exists a $(\lambda_1, \lambda_2) \in \ell^2(X)$ such that

$$P = ([\mathcal{F}_X^*]_1 \mathbf{e}^+ + [\mathcal{F}_X^*]_2 \mathbf{e}^-) (\lambda_1 \mathbf{e}^+ + \lambda_2 \mathbf{e}^-)$$

$$= [\mathcal{F}_X^*]_1 \lambda_1 \mathbf{e}^+ + [\mathcal{F}_X^*]_2 \lambda_2 \mathbf{e}^-.$$

Then c is unique minimizer to the problem (2).

Lemma

If $N \geq |T|$ then \mathcal{F}_{TX} is injective almost surely.

Moreover, we introduce the **restriction** operator
 $R_T : \ell^2([-q, q]^2) \rightarrow \ell^2(T)$, $R_T c_{k\tilde{k}} = c_{k\tilde{k}}$ for $(k, \tilde{k}) \in T$.

Its adjoint $R_T^* = E_T : \ell^2(T) \rightarrow \ell^2([-q, q]^2)$ is the operator that extends a vector outside T by zero, i.e., $(E_T d)_{k\tilde{k}} = d_{k\tilde{k}}$ for $(k, \tilde{k}) \in T$ and $(E_T d)_{k\tilde{k}} = 0$ otherwise.

Assuming that $\mathcal{F}_{TX}^* \mathcal{F}_{TX} : \ell^2(T) \rightarrow \ell^2(T)$ is invertible, we construct, P explicitly by

$$P := \mathcal{F}_X^* \mathcal{F}_{TX} (\mathcal{F}_{TX}^* \mathcal{F}_{TX})^{-1} R_T \text{sgn}(c), \text{ with}$$

$$\mathcal{F}_X^* \mathcal{F}_{TX} := \mathbf{e}^+ [\mathcal{F}_X^* \mathcal{F}_{TX}]_1 + \mathbf{e}^- [\mathcal{F}_X^* \mathcal{F}_{TX}]_2 \text{ and}$$

$$\mathcal{F}_{TX}^* \mathcal{F}_{TX} := \mathbf{e}^+ [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_1 + \mathbf{e}^- [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_2$$

Auxiliary operators and expectation value

We start with some **auxiliary operators**:

$$H := \ell^2(T) \rightarrow \ell^2([-q, q]^2) \text{ and } H_0 := \ell^2(T) \rightarrow \ell^2(T),$$

$$I := \mathbf{e}^+ I + \mathbf{e}^- I \text{ and } E := \mathbf{e}^+ E + \mathbf{e}^- E.$$

$$H_0 = \mathbf{e}^+ (NI - [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_1) + \mathbf{e}^- (NI - [\mathcal{F}_{TX}^* \mathcal{F}_{TX}]_2),$$

$$H := \mathbf{e}^+ (NE - [\mathcal{F}_X^* \mathcal{F}_{TX}]_1) + \mathbf{e}^- (NE - [\mathcal{F}_X^* \mathcal{F}_{TX}]_2).$$

The **expectation value** of a bicomplex random variable function are represented as

$$\begin{aligned} \mathbb{E}[f(\zeta_1 + \mathbf{j}\zeta_2)] &:= \text{Sc} \left[\int_{C_1} f_1(z) \frac{dz}{z} \mathbf{e}^+ + \int_{C_2} f_2(w) \frac{dw}{w} \mathbf{e}^- \right] \\ &= \frac{1}{4\pi\mathbf{i}} \int_0^{2\pi} f_1(e^{\mathbf{i}x^1}) dx^1 + \frac{1}{4\pi\mathbf{i}} \int_0^{2\pi} f_2(e^{\mathbf{i}x^2}) dx^2. \end{aligned}$$

Proof - probability results 1

The study of the powers of H leads to

$$\mathbb{E}_X[\|H_0^n\|_F^2] = \sum_{k=1}^{\min\{n,N\}} \frac{N!}{(N-k)!} \sum_{\mathcal{A} \in P(2n,k)} \mathcal{C}(\mathcal{A}, T),$$

where $\mathcal{C}(\mathcal{A}, T)$ denotes

$$\sum_{\substack{k_1, k_2, \dots, k_{2n} \in T \\ k_{r+1}^1 \neq k_r^1, k_{s+1}^2 \neq k_s^2}} \times \prod_{A \subset \mathcal{A}} \left[\delta \left(\sum_{r \in A} (k_{r+1}^1 - k_r^1) \right) + \delta \left(\sum_{s \in A} (k_{s+1}^2 - k_s^2) \right) \right].$$

Proof - probability results

Since $\|(N^{-1}H_0)^n\|_F \leq \kappa$ (< 1) means \mathcal{F}_{TX} is injective

$$\begin{aligned} P(\|(N^{-1}H_0)^n\|_F \geq \kappa) &= P(\|(H_0)^n\|_F^2 \geq N^{2n}\kappa^2) \\ &\leq N^{-2n}\kappa^{-2n}\mathbb{E}_X[\|(H_0)^n\|_F^2] \end{aligned}$$

Using the argument that the **probability of failure** can be bounded by

- the probability of \mathcal{F}_{TX} **not** being injective, and

- $\sup_{(k,\tilde{k}) \in T^c} |P_{k\tilde{k}}| \geq 1$

we get the desired result.

Applications - Bicomplex

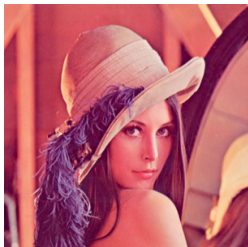


Figure: The original image - $N^2 = 262144$ pixels (512×512)

The computations were made on a computer with Intel(R) Core(TM) i7-4790U CPU 3.60 GHz, RAM 16 GB, Windows 8.1, OS 64-bit(win64), and running Matlab R2020b.

Applications - Bicomplex

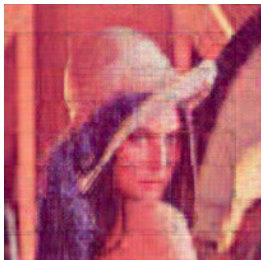


Figure: The reconstructed image with 15% of the total information

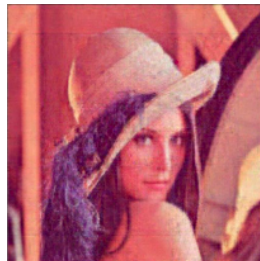


Figure: The reconstructed image with 30% of the total information

Applications - Bicomplex



Figure: The reconstructed image with 50% of the total information

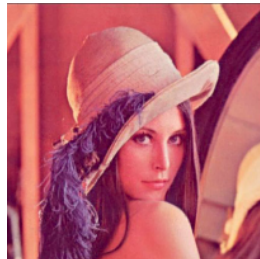


Figure: The reconstructed image with 75% of the total information

Quaternionic Setting

- Compressed Sensing for Quaternionic Signals, Narciso Gomes, Uwe Kähler and Stefan Hartmann, Complex Analysis and Operator Theory. 11, 417-455(2017)
- <https://link.springer.com/article/10.1007/s11785-016-0607-7>

Quaternionic algebra

The quaternionic algebra is an extension of complex numbers to a 4D algebra. Every element of \mathbb{H} is a linear combination of a real and three orthogonal imaginary units (denoted by $\mathbf{i}, \mathbf{j}, \mathbf{k}$) with real coefficients

$$\mathbb{H} = \{q : q = q_0 + \mathbf{i}q_1 + \mathbf{j}q_2 + \mathbf{k}q_3, \quad q_0, q_1, q_2, q_3 \in \mathbb{R}\},$$

where the elements $\mathbf{i}, \mathbf{j}, \mathbf{k}$ obey the Hamilton's multiplication rules

$$\mathbf{i}\mathbf{j} = -\mathbf{j}\mathbf{i} = \mathbf{k}, \mathbf{j}\mathbf{k} = -\mathbf{k}\mathbf{j} = \mathbf{i}, \mathbf{k}\mathbf{i} = -\mathbf{i}\mathbf{k} = \mathbf{j}, \mathbf{i}^2 = \mathbf{j}^2 = \mathbf{k}^2 = \mathbf{i}\mathbf{j}\mathbf{k} = -1.$$

The **vectorial part** is denoted as $\text{Vec}(q) = \mathbf{i}q_1 + \mathbf{j}q_2 + \mathbf{k}q_3$. The **scalar part** is denoted as $\text{Sc}(q) = q_0$.

\mathbb{H} -valued function space

\mathbb{H} -conjugation of a given

$q = q_0 + \mathbf{i}q_1 + \mathbf{j}q_2 + \mathbf{k}q_3 = Sc(q) + Vec(q)$ is

$$\bar{q} = Sc(q) - Vec(q).$$

Consider the quaternion-valued left-Hilbert module $L_2(\mathbb{R}^2; \mathbb{H})$ equipped with the quaternionic-valued inner product

$$(f, g) := \int_{\mathbb{R}^2} f(\mathbf{x}) \overline{g(\mathbf{x})} d\mathbf{x}.$$

Additionally, consider also the real-valued inner product

$$\langle f, g \rangle := Sc(f, g) = \int_{\mathbb{R}^2} Sc[f(\mathbf{x}) \overline{g(\mathbf{x})}] d\mathbf{x}.$$

Our setting and goal

Denote by Π_q the space of functions spanned by trigonometric quaternionic polynomials P of degree at most q .

Now, we consider the signal

$$P(x, y) = \sum_{(k,m) \in T} \left(e^{ikx} e^{jmy} \right) c_{k,m},$$

with the sequence $(c_{k,m})$, $c_{k,m} \in \mathbb{H}$ having support on T and $(x, y) \in \mathbb{R}^2$.

Given the samples $P(x_j, y_j), j = 1, \dots, N$, we want to reconstruct P by means of an ℓ_1 -minimization process.

Main theorem

Theorem

Assume $f \in \prod_q$ with some sparsity $M \in \mathbb{N}$, and consider $X := \{x_1, x_2, \dots, x_N\}$ a sampling set of *l.i.r.v.* with uniform distribution on $[0, 2\pi]^2$.

Choose $n \in \mathbb{N}$, $\beta > 0$, $\kappa > 0$ and $K_1, K_2, \dots, K_n \in \mathbb{N}$ s. t.

$$a := \sum_{m=1}^n \beta^{n/K_m} < 1 \quad \text{and} \quad \frac{\kappa}{1 - \kappa} \leq \frac{1 - a}{1 + a} M^{-3/2}.$$

For $\theta := N/M$ we have that f can be reconstructed with probability at least

$$1 - \left(\beta^{-2n} \sum_{m=1}^n G_{2mK_m}(\theta) + M\kappa^{-2} G_{2n}(\theta) \right)$$

from its sample values $f(x_1), \dots, f(x_N)$ by ℓ^1 -minimization.

How do we apply our fundamental lemma?

- Sampling operator

$$\mathcal{F}_X c(x_i, y_i) = \sum_{k, m \in D} e^{ikx_i} e^{jmy_i} c_{k, m}.$$

- Auxiliary operators:

$$\begin{aligned} H &:= \ell^2(T) \rightarrow \ell^2(D) \quad \text{and} \quad H_0 := \ell^2(T) \rightarrow \ell^2(T) \\ H &:= NE_T - \mathcal{F}_X^* \mathcal{F}_{TX} \quad \text{and} \quad H_0 = NI - \mathcal{F}_{TX}^* \mathcal{F}_{TX}, \end{aligned}$$

with

- \mathcal{F}_{TX} the restriction of \mathcal{F}_X supported only on T .
- $(E_T d)_k = d_k$ for $k \in T$ and $(E_T d)_k = 0$ for $k \notin T$.

Proof - probability results 1

The study of the powers of H_0 leads to

$$\mathbb{E}_X[\|H_0^n\|_F^2] = \sum_{k=1}^{\min\{n, N\}} \frac{N!}{(N-k)!} \sum_{\mathcal{A} \in P(2n, k)} \mathcal{C}_{\mathbb{H}}(\mathcal{A}, T),$$

where $\mathcal{C}_{\mathbb{H}}(\mathcal{A}, T)$ is bounded by

$$\mathcal{C}_{\mathbb{H}}(\mathcal{A}, T) \leq |T|^{2n-k+1} \leq M^{2n-k+1}.$$

Proof - probability results

Since $\|(N^{-1}H_0)^n\|_F \leq \kappa$ (< 1) means \mathcal{F}_{TX} is injective

$$\begin{aligned} P(\|(N^{-1}H_0)^n\|_F \geq \kappa) &= P(\|(H_0)^n\|_F^2 \geq N^{2n}\kappa^2) \\ &\leq N^{-2n}\kappa^{-2n}\mathbb{E}_X[\|(H_0)^n\|_F^2] \end{aligned}$$

Using the argument that the **probability of failure** can be bounded by

- the probability of \mathcal{F}_{TX} **not** being injective, and

- $\sup_{(k,\tilde{k}) \in T^c} |P_{k\tilde{k}}| \geq 1$

we get the desired result.

Applications - Quaternions



Figure: The original image

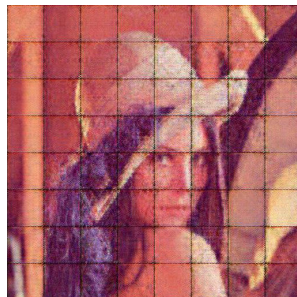


Figure: The reconstructed image.

The original and reconstructed images from Lena with $N^2 = 262144$ pixels (512×512). For the reconstruction we use 40000 pixels ($M = 625$ samples in each block) which corresponds to $\approx 15.26\%$ of the total information.

Applications - Quaternions

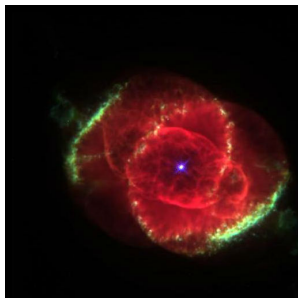


Figure: The original image

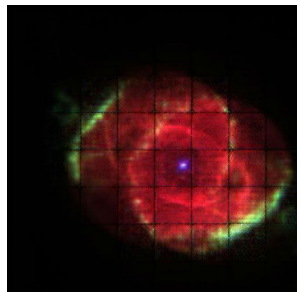


Figure: The reconstructed image.

The original and reconstructed images from Galaxy with $N^2 = 262144$ pixels (512×512). For the reconstruction we use 40000 pixels ($M = 625$ samples in each block) which corresponds to $\approx 15.26\%$ of the total information.

Applications - Quaternions



Figure: The original image

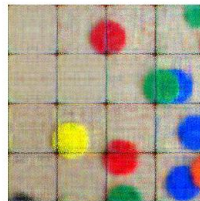


Figure: The reconstructed image.

The original and reconstructed images from Coloured chips with $N^2 = 65536$ pixels (256×256). For the reconstruction we use 10000 pixels ($M = 625$ samples in each block) which corresponds to $\approx 15.26\%$ of the total information.

Conclusion

Further challenges

- Extension to higher dimensions is more difficult since we loose the norm property
 - $\|X^n\| \leq \|X\|^n$
 - $\|XY\| \leq \|X\| \|Y\|$
 - $\|\lambda X\| \leq |\lambda| \|X\|$

Thank you!

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